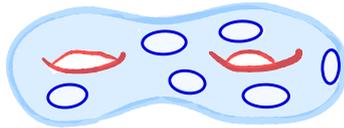


Reminder:

- Knots and links:



- Surfaces

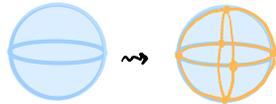


↪ genus $g=2$

↪ # boundary components $b=6$

- Euler characteristic:

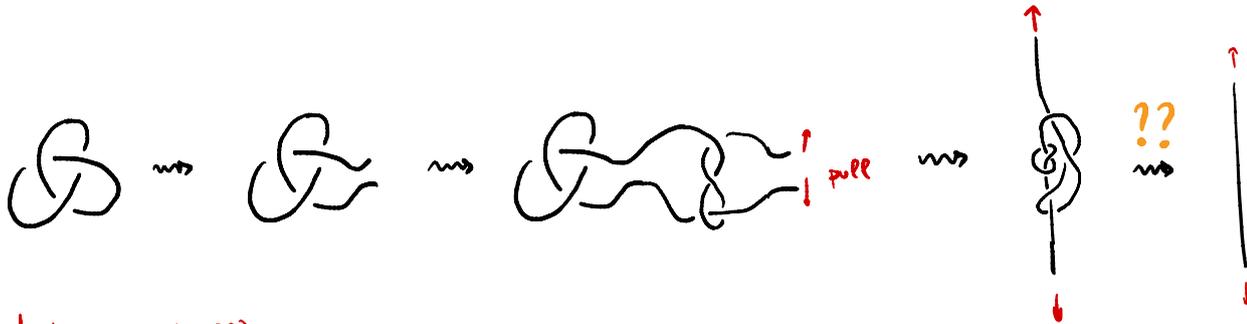
1. Find a triangulation



2. $\chi = V - E + F$

- $\chi(T_{g,b}) = 2 - 2g - b$

New Question: Can you untie a knot by knotting it more?



Intuition (poll)

Any ideas?

Mathematical formulation:



Are there two (proper) knots K_1, K_2 such that

$$K_1 \# K_2 = \text{unknot?}$$

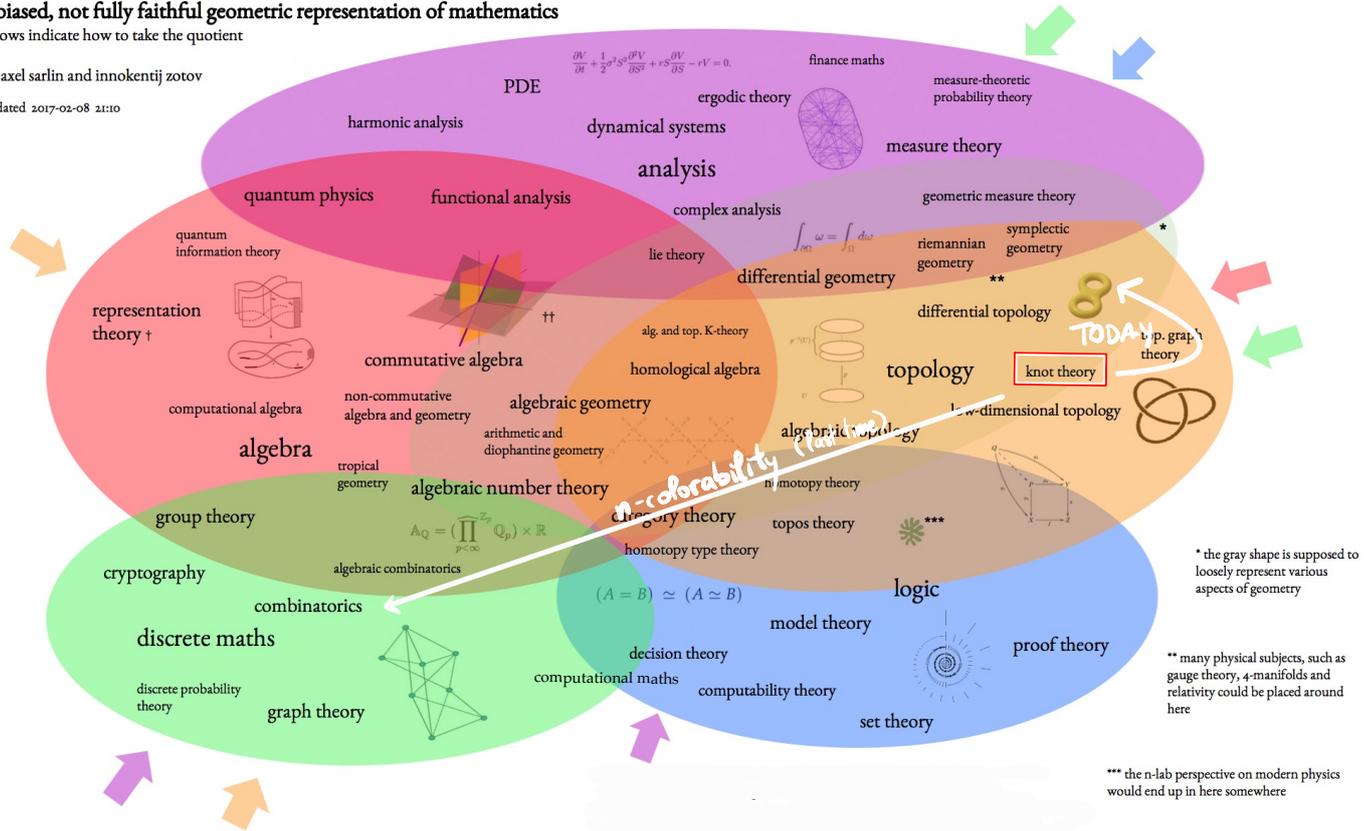
Our tool will come from the topology of surfaces

a biased, not fully faithful geometric representation of mathematics

arrows indicate how to take the quotient

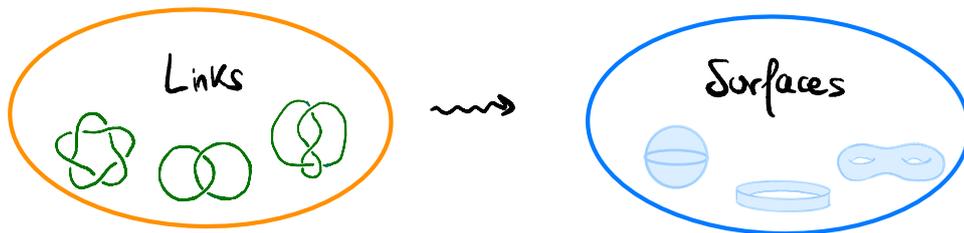
by axel sarlin and innokentij zotov

updated 2017-02-08 21:10



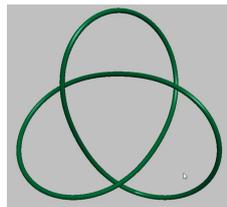
7. Seifert surfaces

Back to our goal:

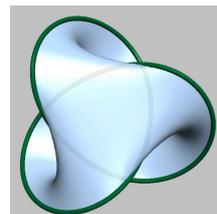


Idea (1930s): [Animation]

Question:



??



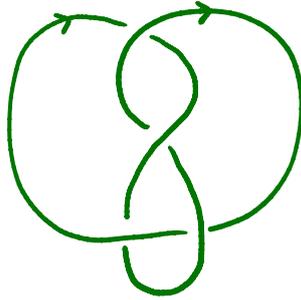
1934, Seifert:



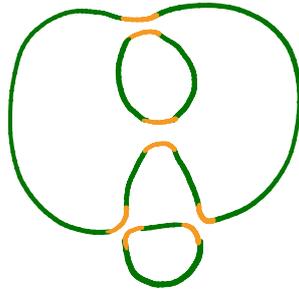
: "Here is an algorithm"

How does it work?

Step 1: Choose a diagram and orientation for the link:



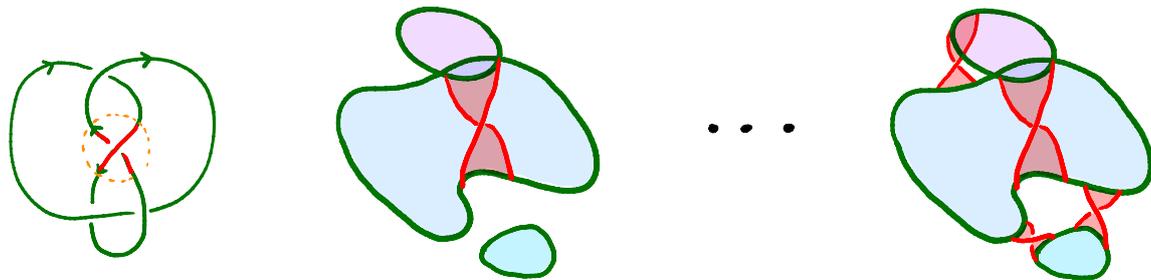
Step 2: smooth out every crossing:  and  become 



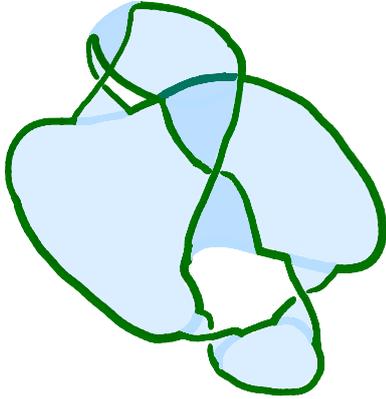
Step 3: make every resultant circle into a disk, and place the disks at different heights



Step 4: place twisted bands where there used to be crossings:



Result:



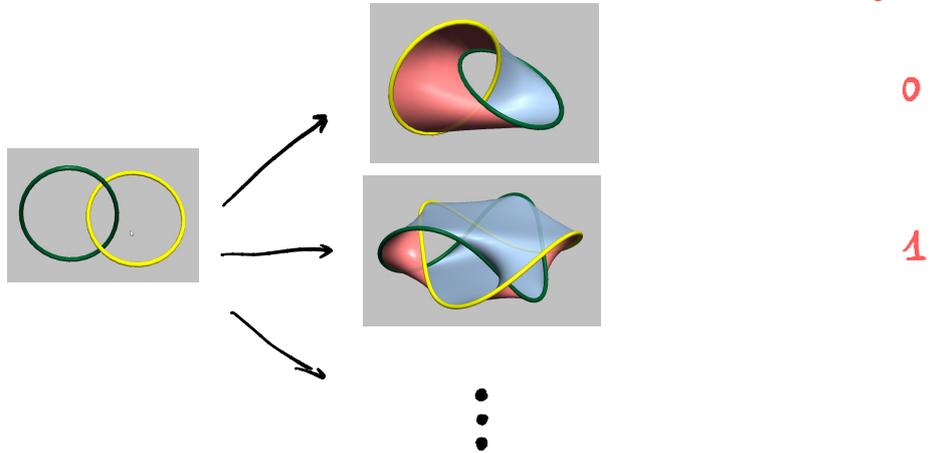
Remark: the result is always **orientable**.

Definition: a surface whose boundary components form a link L is called a **Seifert surface**.

Example: the Seifert algorithm gives a Seifert surface

Warning: These are not unique!

The theory of surfaces gives us a new invariant:



Definition: the **genus** of a link $g(L)$ is the **minimum** genus of a Seifert surface for L .

[Seifert surfaces software]

Examples:

• Unknot:  =  = $T_{0,1}$ 0 is obviously minimum.

• Trefoil:  \rightsquigarrow  = $T_{1,1}$ so $g(\text{trefoil}) \leq 1$.

Observation: If $g(K) = 0$, then $K = \text{unknot}$. Proof: $g(K) = 0$ means that K is the boundary of $T_{0,1} =$ . But  is clearly unknotted.

Consequence: $g(\text{trefoil}) = 1$.

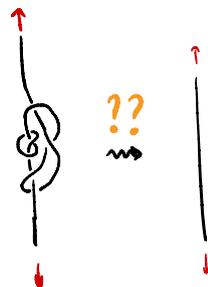
Remark: Sage gives the genus of a link: `L.genus()`

First task: Euler characteristic of Seifert surfaces.

8. Applications: no-unknotting and knot primality

Recall our original question:

Can I untie a knot by knotting it more?



Crucial Theorem: $g(K_1 \# K_2) = g(K_1) + g(K_2)$

Corollary: If K_1 and K_2 are knots but $K_1, K_2 \neq \text{unknot}$, then $K_1 \# K_2 \neq \text{unknot}$.

Proof: $K_1 \neq \bigcirc \Rightarrow g(K_1) \geq 1$.

$K_2 \neq \bigcirc \Rightarrow g(K_2) \geq 1$.

Therefore $g(K_1 \# K_2) \stackrel{\text{Thm}}{=} g(K_1) + g(K_2) \geq 2 > 0 = g(\bigcirc)$ \square

Definition: A knot K is prime if whenever $K = K_1 \# K_2$, then either $K_1 = \text{unknot}$ or $K_2 = \text{unknot}$.

Examples:  $\neq K_1 \# K_2$ prime

 $=$  $\#$  not prime

TABLE OF PRIME KNOTS UP TO 8 CROSSINGS

							
0 ₁	3 ₁	4 ₂	5 ₁	5 ₂	6 ₁	6 ₂	6 ₃
							
7 ₁	7 ₂	7 ₃	7 ₄	7 ₅	7 ₆	7 ₇	7 ₈
							
8 ₁	8 ₂	8 ₃	8 ₄	8 ₅	8 ₆	8 ₇	8 ₈
							
8 ₉	8 ₁₀	8 ₁₁	8 ₁₂	8 ₁₃	8 ₁₄	8 ₁₅	8 ₁₆
							
8 ₁₇	8 ₁₈	8 ₁₉	8 ₂₀	8 ₂₁			

Remark: Sage has the prime knots stored up to ~ 14 crossings.

No one has classified the primes with ≥ 17 crossings.

Crucial Theorem: $g(K_1 \# K_2) = g(K_1) + g(K_2)$

Corollary: Every knot can be written as a sum of prime knots.

Proof: Let K be a knot. If $g(K) = 1$, then K is prime (you prove this in the exercises)

This proves the claim for knots of genus 1.

If $g(K) = 2$, then either $K = K_1 \# K_2$ for $K_1, K_2 \neq \text{unknot}$, in which case \downarrow $g(K_1) = g(K_2) = 1$ and each of K_1, K_2 are prime. Otherwise K itself is prime.
 by the theorem

This proves the claim for knots of genus 2.

If $g(K) = 3$, then either $K = K_1 \# K_2$ for $K_1, K_2 \neq \text{unknot}$, in which case $g(K_1), g(K_2) \leq 2$ and each of K_1, K_2 are prime. Otherwise K itself is prime.

... \square

This is called a "proof by induction".

Remark: Seifert surfaces can be used to show uniqueness.

Proof of the crucial theorem

Warning: this is a step harder than what we've done so far, it's supposed to give you a taste!

Definition: Let S be a surface with a disk D bounding a circle in S . Then "performing surgery along D " is the action of replacing S by a surface S' where we cut the surface S at the circle and cap off the remaining "tubes":



Lemma: If performing surgery on a surface S results in a connected surface S' , then $g(S') = g(S) - 1$.

Proof of the lemma: Take a triangulation for S s.t. the disk is one of the triangles. Then after surgery:

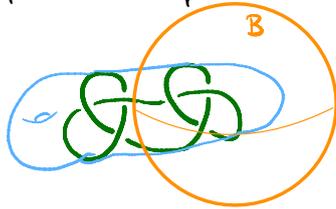
$V \rightsquigarrow V+3$, $E \rightsquigarrow E+3$, $F \rightsquigarrow F+2$. So $\chi \rightsquigarrow \chi+2$. But $\chi = 2-2g$ so $g \rightsquigarrow g-1$. \square

Proof of the theorem: It suffices to show $g(K_1 \# K_2) \geq g(K_1) + g(K_2)$.

Take a Seifert surface S for $K_1 \# K_2$ of minimal genus:



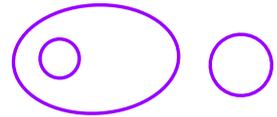
Next, take a sphere that separates the two summed knots:



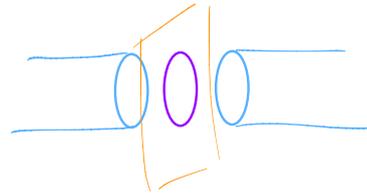
The intersection between B and S consists of:



+ some circles

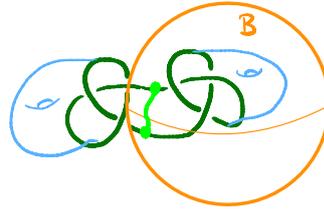


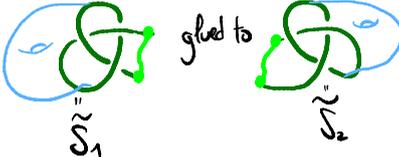
Take an innermost circle. Then perform surgery along it:



The resulting surface is a Seifert surface for $K_1 \# K_2$. If it is connected, then by the Lemma, its genus is lower than $g(S)$, impossible. It follows that it is disconnected and therefore only one of the components bounds $K_1 \# K_2$, and we can remove the other one.

This procedure removes innermost circles one at a time, so eventually we get a surface \tilde{S} whose intersection with B is just  :



Next, notice that $\tilde{S} = \tilde{S}_1 \cup \tilde{S}_2$ glued to 

Now $g(K_1 \# K_2) = g(\tilde{S}) = g(\tilde{S}_1) + g(\tilde{S}_2) \geq g(K_1) + g(K_2)$ as desired. \square

- Your second task:
- Investigate the genus of knots
 - Prove the infinitude of prime knots.