

§3 Chebyshev's Estimates for $\pi(x)$

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Half way between Gauss' conjecture and the proof of the Prime Number Theorem, Chebyshev was able to prove a weaker version of PNT. This was one of his biggest achievements, and, for 50 years, the most that was known about $\pi(x)$.

JHM (Chebyshev, 1850) There are positive constants c & C so that

$$c \frac{x}{\log x} \leq \pi(x) \leq C \frac{x}{\log x}, \text{ for } x \geq 2.$$

The proof that $\pi(x) \rightarrow \infty$ took two lines. This takes four pages.

LEMMA (About the middle binomial coefficient). For all $N \in \mathbb{N}$,

$$(a) \quad \frac{2^{2N}}{2N} \leq \binom{2N}{N} \leq 2^{2N};$$

$$(b) \quad N^{\pi(2N) - \pi(N)} \leq \binom{2N}{N} \leq (2N)^{\pi(2N)}.$$

Proof (a). According to the Binomial Theorem, for $n \in \mathbb{N}$ and $x \in \mathbb{R}$,

$$(1+x)^n = \sum_{k=0}^n \binom{n}{k} x^k, \quad \text{with } \binom{n}{k} = \frac{n!}{k!(n-k)!}.$$

Putting $n = 2N$ and $x = 1$, this gives

$$\checkmark \quad 2^{2N} = \sum_{k=0}^{2N} \binom{2N}{k},$$

from which the right inequality in (a) follows.

Observe that

$$\binom{n}{k-1} \leq \binom{n}{k} = \binom{n}{n-k} \text{ for } 0 < k \leq \frac{n}{2}.$$

It follows that the middle term in \checkmark is largest and is ≥ 2 (= sum of terms with $k=0$ & $k=2N$). Thus the sum in \checkmark is $\leq 2N \binom{2N}{N}$, proving the left inequality in (a).

EXERCISE 1. Prove the observation $\textcircled{<}$, using the formulas for $\binom{n}{k}$ and $\binom{n}{k-1}$.

(b) Since $(n!)^2 \binom{2N}{N} = (2N)!$, and $n!$ is divisible by each prime $p \leq n$ and no prime $p > n$, it follows that $\binom{2N}{N}$ is divisible by each prime p with $N < p \leq 2N$, and therefore by their product, so

$$\binom{2N}{N} \geq \prod_{N < p \leq 2N} p \geq N^{\pi(2N) - \pi(N)},$$

giving the left inequality in (b).

To prove the right inequality in (b) we first get the prime factorization of $n!$:

$$n! = \prod p^{\sigma_p(n)}, \quad \text{with } \sigma_p(n) = \sum_{r \in \mathbb{N}} \left[\frac{n}{p^r} \right].$$

To see this, observe first that the multiplicity of a prime p in the prime factorization of a positive integer k equals the number of positive integers r with $p^r | k$.

It follows that the multiplicity of p in $n!$ equals the number of pairs of positive integers r, k for which $p^r | k$ and $k \leq n$. For fixed r , there are $\left[\frac{n}{p^r} \right]$ positive integers k with $p^r | k$ and $k \leq n$.

It follows from \ast that

$$n! = \prod p^{\tau_p(n)}, \quad \text{with } \tau_p(n) = \sum_{r \in \mathbb{N}} \left(\left[\frac{2N}{p^r} \right] - 2 \left[\frac{N}{p^r} \right] \right).$$

Observe next that the function θ defined on \mathbb{R} by

$$\theta(u) = [2u] - 2[u]$$

satisfies $\theta(u+1) = \theta(u)$ (since $[v+1] = [v] + 1$), and $\theta(u) = 0$ on $[0, \frac{1}{2})$, 1 on $[\frac{1}{2}, 1)$.

So $0 \leq \theta(u) \leq 1$ for all $u \in \mathbb{R}$. Since only terms with $p^r \leq 2N$ can contribute to $\tau_p(n)$,

we get $\tau_p(n) \leq \left[\frac{\log 2N}{\log p} \right]$. Therefore, from $\ast\ast$,

$$\binom{2N}{N} \leq \prod_{p \leq 2N} p^{\frac{\log 2N}{\log p}} = \Theta(2N)^{\pi(2N)},$$

giving the right inequality in (b).

EXERCISE 2. Prove the \Leftarrow .

Proof of Chebyshev's Theorem: The four inequalities in the lemma can be combined, in two different ways, to give information ($*$ & $**$ below) about $\pi(x)$.

Combining the left inequality in (a) with the right inequality in (b) gives

$$\frac{2^{2N}}{2N} \leq (2N)^{\pi(2N)}$$

Taking logs (& using $\log u \leq \log v$ for $u \leq v$), we get

$$\pi(2N) \geq (\log 2) \frac{2N}{\log 2N} - 1, \text{ (so)}$$

$*$ $\left[\pi(2N) \geq b \frac{2N}{\log 2N}, \text{ for all } N \in \mathbb{N} \right]$

with $b = \frac{\log 2}{2}$, since $\frac{\log 2}{2} \geq \frac{\log 2N}{2N}$ for $N \in \mathbb{N}$.

EXERCISE 3. Verify \Leftarrow by showing that $\frac{\log x}{x}$ decreases for $x \geq e$ and $\frac{\log 2}{2} = \frac{\log 4}{4}$.

Then show how to use \uparrow to derive $*$ from the inequality just above $*$, just before (so).

Combining the left inequality in (b) with the right inequality in (a) gives

$$N^{\pi(2N) - \pi(N)} \leq 2^{2N}$$

Taking logs,

$$\pi(2N) - \pi(N) \leq (\log 4) \frac{N}{\log N}$$

We will use this to show that there is a positive constant B so that

$**$ $\left[\pi(2^k) \leq B \frac{2^k}{\log 2^k}, \text{ for all } k \in \mathbb{N} \right]$

For $k=1$ or 2 this is easy, with $B = \log 2$. We try induction. For $k \geq 2$, the inequality just above $**$ with $N = 2^{k-1}$ gives

$$\pi(2^{k+1}) = \pi(2^{k+1}) - \pi(2^k) + \pi(2^k)$$

$$\log 2^k \text{ means } \log(2^k) \leq (\log 4) \frac{2^k}{\log 2^k} + B \frac{2^k}{\log 2^k} \stackrel{?}{\leq} B \frac{2^{k+1}}{\log 2^{k+1}}$$

Here we have supposed that $**$ is true for k , with some B , and ask if it is true for $k+1$, with the same B . Cancelling $\frac{2^k}{\log 2}$ from the questioned last inequality gives

$$\frac{\log 4}{k} + \frac{B}{k} \stackrel{?}{\leq} \frac{2B}{k+1}$$

Multiplying by $k(k+1)$ gives

$$(\log 4)(k+1) \stackrel{?}{\leq} (k-1)B,$$

i.e.

$$B \stackrel{?}{\geq} (\log 4) \frac{k+1}{k-1}.$$

For $k \geq 2$ the right side is largest for $k=2$, so, yes, it works, for $B = 3 \log 4$.

Finally, we use $*$ and $**$ to finish the proof of Chebyshev's Theorem.

For $x \geq 2$, choose $N \in \mathbb{N}$ with $2N \leq x \leq 4N$. Then by $*$,

$$\pi(x) \geq \pi(2N) \geq b \frac{2N}{\log 2N} = \frac{b}{2} \frac{4N}{\log 2N} \geq c \frac{x}{\log x},$$

with $c = \frac{b}{2} (= \frac{\log 2}{4} > 0.17)$, giving the left inequality in the theorem.

For $x \geq 2$, choose $k \in \mathbb{N}$ with $2^k \leq x \leq 2^{k+1}$. Then by $**$,

$$\pi(x) \leq \pi(2^{k+1}) \leq B \frac{2^{k+1}}{\log 2^{k+1}} = 2B \frac{2^k}{\log 2^{k+1}} \leq C \frac{x}{\log x},$$

with $C = 2B (= 12 \log 2 < 8.4)$.

For more on $\pi(x)$, see H. Davenport's Multiplicative Number Theory.