

## NOTES ON THE POHLIG-HELLMAN ATTACK ON THE DISCRETE LOG PROBLEM

Let  $p =$  large prime and  $1 < \alpha < p$  a primitive root (mod  $p$ ). If  $\alpha, \beta, p$  are known and

$$\alpha^x \equiv \beta \pmod{p},$$

the Discrete Log Problem (DLP) is to find  $x$  with  $1 < x < p$ .

**Solving DLP with a brute force attack of  $R$  tries:** *Let's assume we know the solution  $x$  is not too big. For example assume we know  $1 \leq x < R$  for some small integer  $R$ . Then we simply compute*

$$\alpha^0 \equiv 1 \pmod{p}, \alpha \pmod{p}, \alpha^2 \pmod{p}, \alpha^3 \pmod{p}, \dots, \alpha^{R-1} \pmod{p}.$$

*One of the above  $R$  numbers has to equal to  $\beta$ . Then we have found  $x$ .*

**Finding  $x \pmod{q}$  with the Pohlig-Hellman attack:**

**Step 1:** Find a small integer  $q$  where  $q$  divides  $p - 1$ .

**Step 2:** Compute  $A = \alpha^{\frac{p-1}{q}} \pmod{p}$ .

**Step 3:** Compute  $B = \beta^{\frac{p-1}{q}} \pmod{p}$ .

**Step 4:** Solve  $A^y \equiv B \pmod{p}$  with a brute force attack of  $q$  tries. Then  $y \equiv x \pmod{q}$ .

**Example:** Consider the Discrete Log Problem:  $2^x \equiv 17 \pmod{61}$ . Find  $x \pmod{5}$ .

**Step 1:**  $5 \mid (61 - 1)$ .

**Step 2:**  $A \equiv 2^{12} \pmod{61} = 9$ .

**Step 3:**  $B \equiv 17^{12} \pmod{61} = 20$ .

**Step 4:** We make 5 tries in trying to solve  $9^y \equiv 20 \pmod{61}$ :

$$9^0 \equiv 1 \pmod{61}, 9^1 \equiv 9 \pmod{61}, 9^2 \equiv 20 \pmod{61}, 9^3 \equiv 58 \pmod{61}, 9^4 \equiv 34 \pmod{61}.$$

We see that  $y = 2$  is the solution. So  $x \equiv 2 \pmod{5}$ . Actually  $x = 47$  is the solution.

Why does this work? The Pohlig-Hellman attack works because

$$\alpha^x \equiv \beta \pmod{p} \implies \alpha^{x \cdot \frac{p-1}{q}} \equiv \beta^{\frac{p-1}{q}} \pmod{p} \implies \alpha^{y \cdot \frac{p-1}{q}} \equiv \beta^{\frac{p-1}{q}} \pmod{p},$$

where  $y \equiv x \pmod{q}$ .

**Finding  $x$  with the Pohlig-Hellman attack:**

*If one finds  $x \pmod{q}$  for sufficiently many coprime integers  $q$  then one may solve for  $x$  using the Chinese Remainder Theorem.*