

Notes on Pollard's $p - 1$ Attack on RSA

Let $n = pq$ be an RSA key where p, q are primes. Pollard's attack is a method to find the factorization of n . The method will work (with a choice of a large integer B) provided:

- $p - 1$ divides $B!$
- $q - 1$ has a prime factor $> B$.

Step 1: Let $a = 2$ and compute $b \equiv a^{B!} \pmod{n}$.

Step 2: Calculate $\gcd(b - 1, n)$. This should give the prime factor p .

Step 3: If p is not found in step 2 repeat steps 1,2 with $a = 3$. If $a = 3$ fails keep trying other values for a .

Why does this work? Pollard's attack works because of Fermat's Little Theorem.

In fact, since $(p - 1) \mid B!$ this implies that $B! = k \cdot (p - 1)$ for some integer k . Recall that $b \equiv 2^{B!} \pmod{pq}$. It follows that

$$b \equiv 2^{B!} \pmod{p} \equiv 2^{k \cdot (p-1)} \equiv 1 \pmod{p}.$$

Hence p divides $b - 1$.

Warning: The method will work only if $b > 1$.

It will certainly be the case that $b > 1$ if 2 is a primitive root \pmod{q} , since we assumed that $\phi(q) = q - 1$ has a prime factor $> B$, which implies $2^{B!} \not\equiv 1 \pmod{q}$.

If $a = 2$ doesn't work then try $a = 3$, etc, until a good choice of a is obtained. Eventually the factorization of n will be found.

EXAMPLE: Factor $n = 2573$ with Pollard's attack and $B = 5$.

Step 1: $B! = 5! = 120 = 64 + 32 + 16 + 8 = 2^6 + 2^5 + 2^4 + 2^3$.

$$2^{2^1} \pmod{2573} \equiv 4$$

$$2^{2^2} \pmod{2573} \equiv 16$$

$$2^{2^3} \pmod{2573} \equiv 256$$

$$2^{2^4} \pmod{2573} \equiv 1211$$

$$2^{2^5} \pmod{2573} \equiv 2484$$

$$2^{2^6} \pmod{2573} \equiv 202$$

$$\begin{aligned} b &\equiv 2^{120} \pmod{2573} \equiv 2^{2^6} \cdot 2^{2^5} \cdot 2^{2^4} \cdot 2^{2^3} \pmod{2573} \\ &\equiv 202 \cdot 2484 \cdot 1211 \cdot 256 \pmod{2573} \\ &\equiv 280 \end{aligned}$$

Step 2: $\gcd(279, 2573) = 31$.

$$\boxed{n = 31 \cdot 83.}$$